

Topic P : KMP , Automata

What we know  
about strings:

- Rolling hash

- Trie

Good for 90%  
of the problems!

What we will introduce  
today:

- KMP

- AC Automata

Review : Tip of your tongue (NAQ '23).



# KMP

## Motivation

Word search :

Find 'Bob' in 'Alice Bob --'

Naive algo :  $O(|w| \cdot |S|)$

Rolling hash : probabilistic\*  $O(|S|)$

Goal : deterministic  $O(|S|)$

(itself not as important,

but the research motivates other probs)

KMP: deterministic  $O(|S|)$

Intuition: Word search is avg. fast  
'Bob' in 'Alice Bob' -  
B only appeared once!

Counterexample: ' $a$ '  $\times 500$  in ' $a$ '  $\times 1000$   
vs.

' $a$ '  $\times 500$  in ' $a$ '  $\times 499 \rightarrow b$

We only need to optimize for bad scenarios

Optimization 1: Throw away everything after a mismatch.

$$w = \text{aaa} \quad S = \text{aabaaa}$$

aaa       $\begin{matrix} a \\ a \\ \cancel{b} \\ a \\ a \end{matrix}$

When we encounter  $b$ , we can throw away  
the matched prefix and start after  $b$ !

Complexity  $O(|S|)$ ?

Optimization 1: Throw away everything after a mismatch.

$$w = \underline{aaa} \quad S = aab \underline{aaa}$$

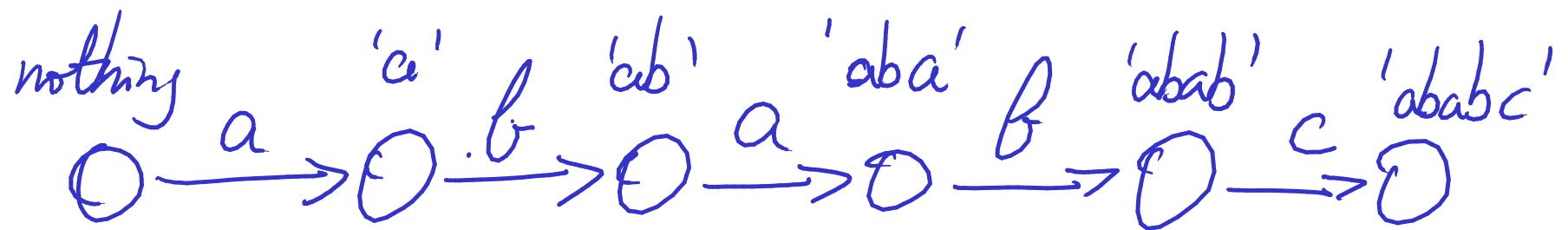
aaa      aab  $\frac{\beta}{\Delta}$  aaa

When we encounter  $\beta$ , we can throw away  
the matched prefix and start after  $\beta$ !

Complexity  $O(|S|)$ ?

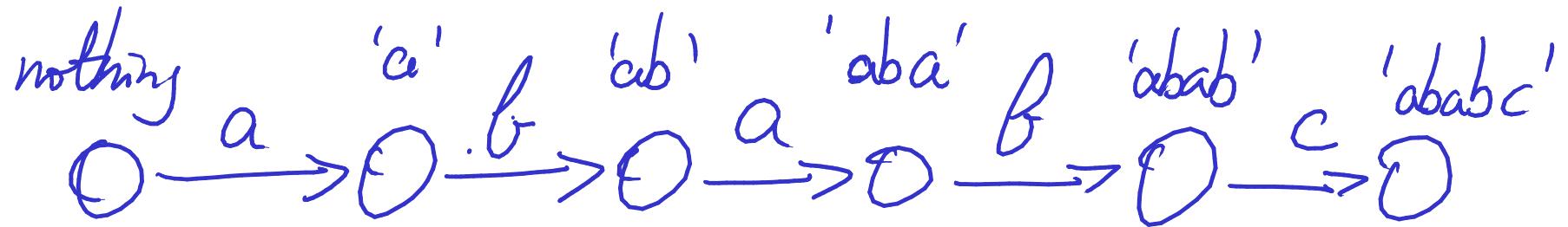
Counterexample 'ababc' in  
'abababc'

## Optimization 2 : Automata - based approach



$$S = abababc$$

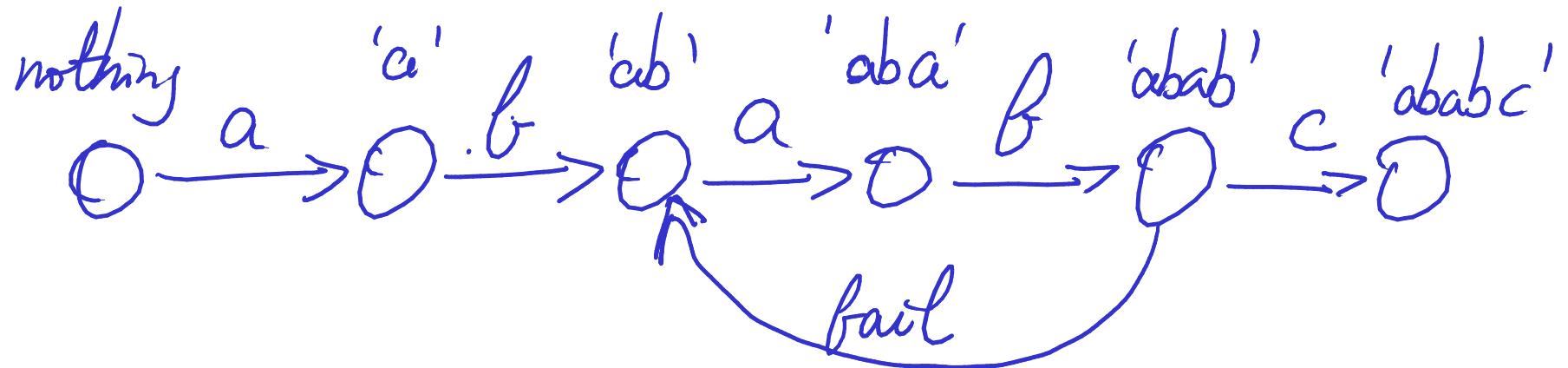
## Optimization 2: Automata-based approach



Problem: need to figure out where to go after failing to advance.

Intuition: go to a 'weaker' node and try to advance again.

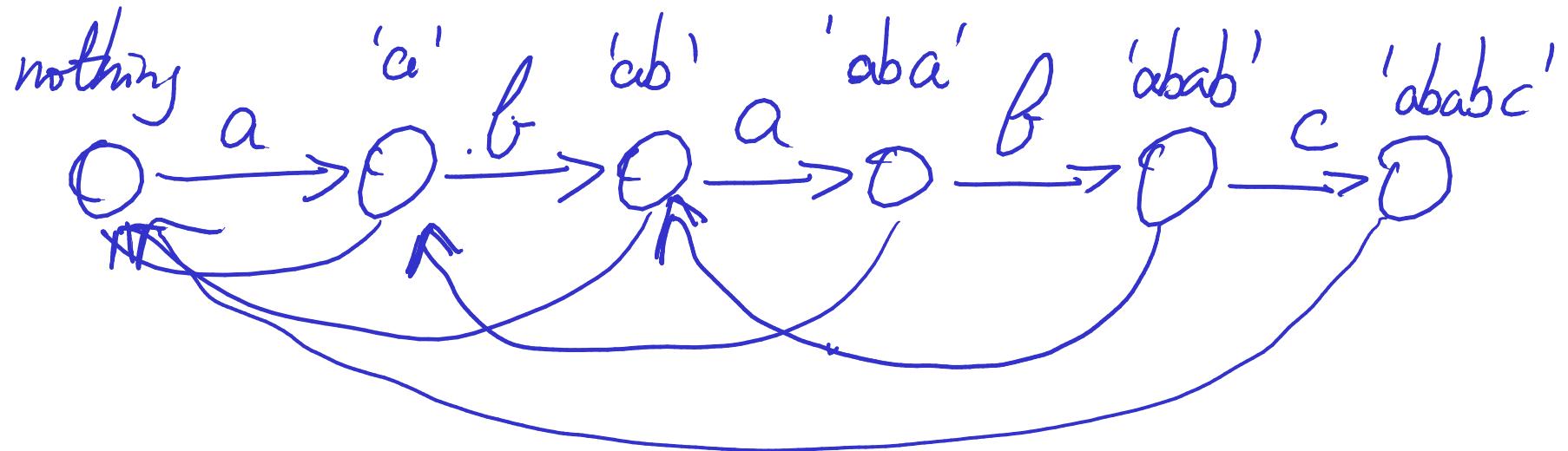
## Optimization 2: Automata-based approach



\* whatever ends with 'abab'  
also ends with 'ab'.

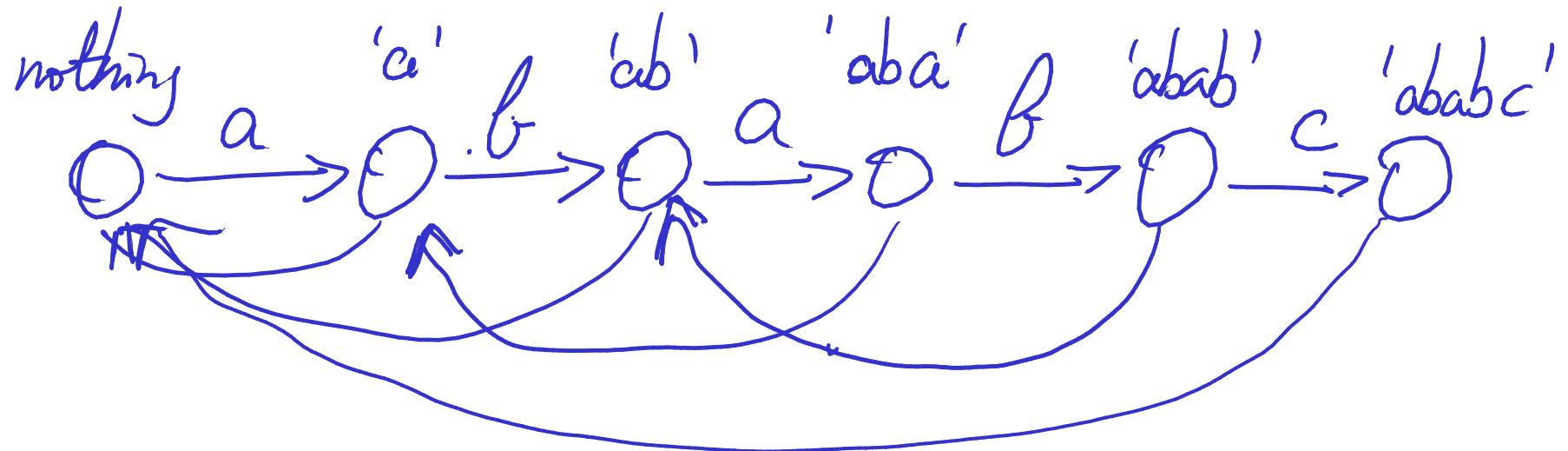
Intuition: go to a 'weaker' node  
and try to advance again.

## Optimization 2 : Automata - based approach



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and try to advance again.

## Optimization 2 : Automata - based approach

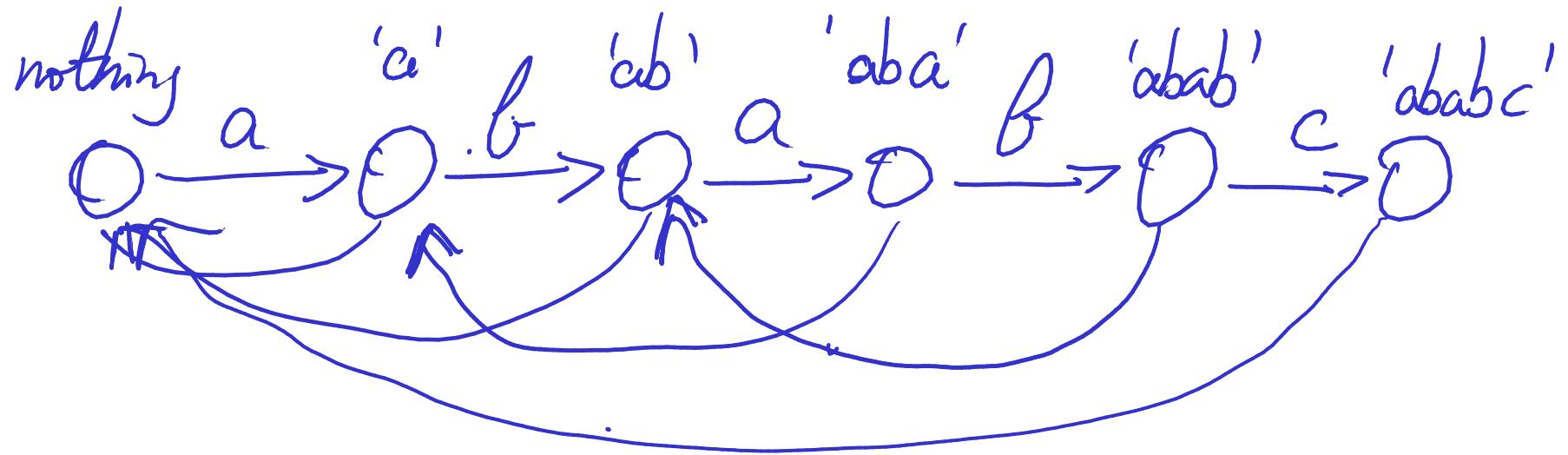


★ fail[i] can be inferred from fail[i-1].

$$'aba' \longrightarrow 'a'$$

$$'aba' + 'b' \longrightarrow 'a' + 'b'$$

## Optimization 2 : Automata - based approach



Refer to  
code base  
for an exact  
code !!!

```
f = fail[i-1], ch = w[i]
while (f is not start) {
    if (f advances with ch) {
        break;
    } f = fail[f];
    if (f advances with ch) {
        fail[i] = f+1;
    } else fail[i] = start;
```

Compress Woods from CF 1200E

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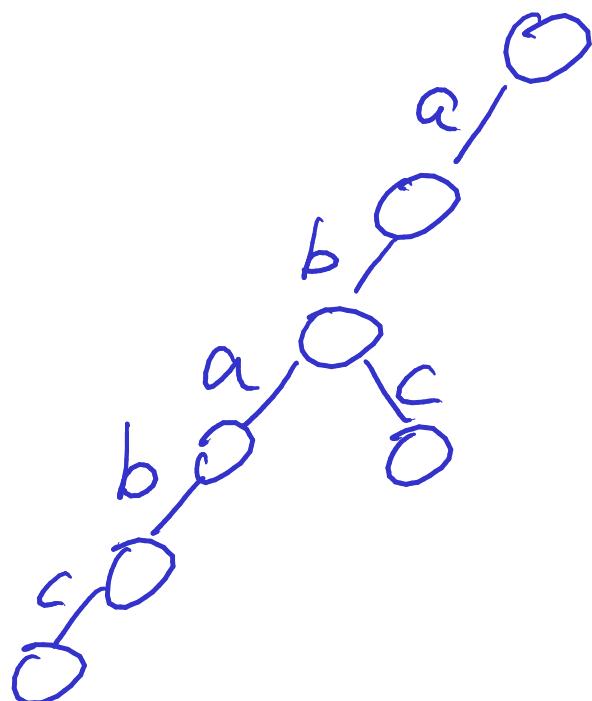


# AC Automata

What if we want to match multiple strings simultaneously?

So, we build on a tree!

e.g. 'ababc' and 'abc'



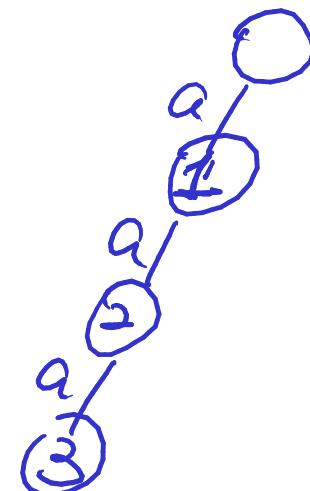
# AC Automata

Caveat 1: a state may match multiple strings.

$$\omega = \{ 'a' \quad 'aa' \quad 'aaa' \}$$

$$S = 'aaa'$$

```
match-ctl[i] =  
    match-ctl[fail[i]] +  
    [i ∈ ω];
```



Caveat 2: Jump to fail is slow in worst.  
Preprocess every jump as gol[i]['a'], etc.

Indie Album on CF 1207G

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